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TID	ITEMS
1	M , N , P , Q
2	M , O
3	M , N , O , Q
4	N , O , Q
5	M , N , O , Q

Table 1: Sample database

III. OBJECTIVE OF THIS PAPER

The main objective of this thesis is to develop a new algorithm called "Binary Enhanced APriori - BEAP Algorithm". ” by combining binary representation and finding rare items from large databases. The secondary objectives are explained below,

- 1) To make a detailed comparison with existing algorithms in terms of running time and memoryconsumption using some standardized databases.
- 2) To analyze and study the current algorithms dealing with rare items mining using Aprioriapproach.

IV. PROBLEM BACKGROUND

Although frequent item mining has evolved to a great extent, finding frequent and rare items from a database is still a complex task because it deals with many calculations and functions related to support counting and pruning. Non-promised goods. The primary problem is generating a large number of frequent items, and this makes the algorithm inefficient. If the minimum support value is set to a very lowvalue, the amount of candidates generated will be high and if the minimum support value is set to a high value, rare items will be missed. This problem is fixed by introducing binary representation in the proposed work and starting to find rare items from the combination of unique items, and the proposed algorithm only handles 2U combinations, i.e. U are unique items.

□

V. FIND UNIQUE OBJECT FROM DB

Individual items from the sample database shown in Table 1 are initially found using the FindUniprocedure shown in Figure 2.

PROCEDURE FindUni (Database D)						
ITEM	T1	T2	T3	T4	T5	
M	1	1	1	0	1	

```

INPUT : Database D
OUTPUT : Unique Items
BEGIN : □
1. Load the input database D into the procedure FindUni
2. Initialize an array G[]= $\phi$ 
3. Find the total number of rows in the DB  $\rightarrow$  T
4. For each row  $r_i \in R$ 
5. For each Item  $I_i \in r_i$ 
6. If  $I_i \text{ Not IN } G[]$  do begin
7.  $G[] = I_i$  // store unique items
8. Else
9. Count  $I_i$  value  $I_i$  counter // count unique items
10. Close IF
11. Close For
12. Close For
Return  $G[]$  with count
END

```

Figure 2: Pseudo code of the procedure to find the unique items.

Figure 2 shows the process of finding individual items from the database. The output of the process is shown in the following table 2.

ITEMS	FREQUENCY
M	4
N	4
O	4
P	1
Q	4

Table 2: Unique items discovered

VI. REPRESENT THE DATA IN BINARY

The database is reconstructed in binary form, i.e. if there is an item in the row identifier, that item is marked as "1". A row identifier marked with "0" does not exist. For example, the unique item {M} is in TID's {T1, T2, T3, T5} and is denoted as , the process of reconstructing the database in binary format is shown in the following Figure 3.

PROCEDURE ReconstructDB (database D, Unique G)

INPUT: Database D, Unique items G **OUTPUT:** Reconstructed DB in binary format **BEGIN:**

1. Load the input database D into the procedure
2. Load the Unique G into the procedure
3. Initialize Binary[] = 0
4. For each row $r_i \in R$
5. For each Item $I_i \in G$
6. if [$I_i = r_i$] do begin
7. Mark I_i "1" and store in Binary
8. Else
9. Mark I_i "0" and store in Binary
10. Close IF
11. Close For
12. Close For

Return Binary[]

END

Figure 3: Process of Reconstructing a Pseudo Code DB in Binary Format

Figure 3 shows the process of reconstructing a database in binary format using the following properties, "If the row identifier contains an item X_i then that item $X_i = 1$ else $X_i = 0$ " The following table 3 shows the reconstructed database and this database is a and contains only zeros.

ITEM	T1	T2	T3	T4	T5	COUNT
M	1	1	1	0	1	4
N	1	0	1	1	1	4
O	0	1	1	1	1	4
P	1	0	0	0	0	1
Q	1	0	1	1	1	4

M	1	1	1	0	1
N	1	0	1	1	1
O	0	1	1	1	1
P	1	0	0	0	0
MNOP	0	0	0	0	0

Table 3: Reconstructed database

VII. SORT THE UPPER LEVEL ITEMS

The permutations start from the upper level (ie) {M, N, O, P, Q}, and the frequency or number of these items is calculated from Table 3, reconstructed using simple and function. For these binary values, as shown in the following section,

The items {{M, N, O, P, Q} are rare items with 0 support. Now the next level items are sorted and the calculations are shown in the following section,

M	1	1	1	0	1
N	1	0	1	1	1
O	0	1	1	1	1
Q	1	0	1	1	1
MNOQ	0	0	1	0	1

The 4-item {M, N, O, P}, {M, N, O, Q}, {M, N, P, Q}, {M, O, P, Q}, {N, O, P, Q} and the count for these items is calculated using binary AND.

B	1	1	1	0	1	
N	1	0	1	1	1	
O	0	1	1	1	1	
P	1	0	0	0	0	
Q	1	0	1	1	1	AND
MNOPQ	0	0	0	0	0	

The set of items {M, N, O, P} is a set of rare items with 0 support. The next item is calculated.

Itemset {M, N, O, Q} = 2 rare items with support. It is a non-zero rare item. The next set of items is calculated.

The itemset {M, N, P, Q} = rare items with support 1 and is non-zero rare items.

M	1	1	1	0	1
N	1	0	1	1	1
P	1	0	0	0	0
Q	1	0	1	1	1
MNPQ	1	0	0	0	0

M	1	1	1	0	1
O	0	1	1	1	1
P	1	0	0	0	0
Q	1	0	1	1	1
MOPQ	0	0	0	0	0

M	1	1	1	0	1
N	1	0	1	1	1
O	0	1	1	1	1
MNO	0	0	1	0	1

M	1	1	1	0	1
N	1	0	1	1	1
P	1	0	0	0	0
MNO	1	0	0	0	0

The itemset {M, O, P, Q} = 0 rare items with support and is zero rare items.

The itemset {N, O, P, Q} = 0 rare items with support and is zero rare items. Now the next level 3-itempermutation is calculated as shown in this section.

3-Items {M, N, O}, {M, N,P}, {M, N, Q}, {M,O, P}, {M, O, Q}, {M,P, Q} , {N,O,P},{N,O,Q}, {N,P,Q},{O,P,Q}.

First {M, N, O} support is calculated as shown below,

N	1	0	1	1	1
O	0	1	1	1	1
P	1	0	0	0	0
Q	1	0	1	1	1
NOPQ	0	0	0	0	0

Items {M, N, O} are rare items with support =2. It is a non-zero rare item. The next set of items iscalculated.

Itemset {M, N, P} = rare items with support 1. It is a non-zero rare item. Similarly all permutations areobtained and calculated. The final results are shown in Table 4.

ITEMS	COUNT	ITEMS	COUNT
MNOPQ	0	MPQ	1
MNOP	0	NOP	0
MNOQ	2	NPQ	1
MNPQ	1	OPQ	0
MOPQ	0	MP	1
NOPQ	0	NP	1
MNO	2	OP	0
MNP	1	PQ	1
MOP	0	P	1
MOQ	2		

Table 4: Final output of rare itemset

VIII. EXPERIMENTAL RESULTS

The entire experiment is executed on a dual core processor with 4GB RAM at 2GHz and the datasets are collected from the following domain <http://fimi.cs.helsinki.fi/data/>. The proposed algorithm is developed using Java SMPF tool and compared with other algorithms like Apriori-Inverse [4] and ARIMA[5] and the results regarding runtime and memory consumption are shown in the following graphs. A synthetic database is created by the IBM-Quest tool and algorithms areexecuted on the synthetic database.

The characteristics of the ranked database are shown in the following Table 5, and these databases are used to implement the algorithms.

Dataset	Size (MB)	# Items	# Trans	AvgTrans
Mushroom	19.3	119	8124	23
Retail	4.2	16,470	88,126	10.3
Pumsb	16.3	2,113	49046	74
Kosarak	30.5	41,271	990,002	8.1

Table 5: The characteristic of the database

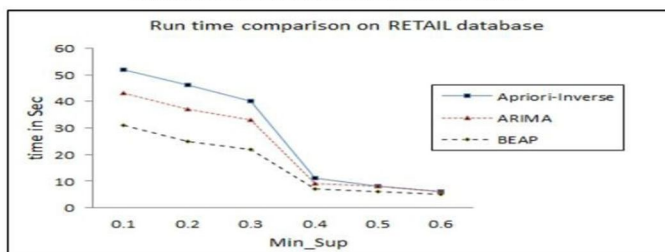


Figure 4: Comparison graph for run time on retail database

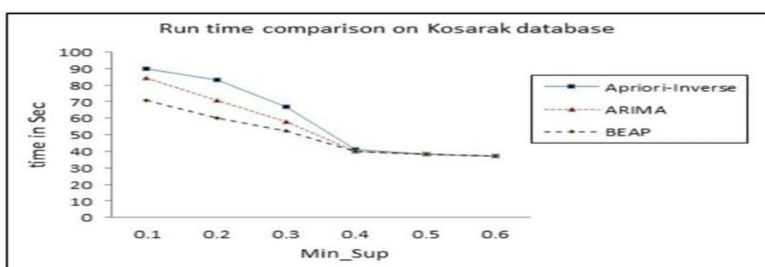


Figure 5: Comparison graph for run time on retail database

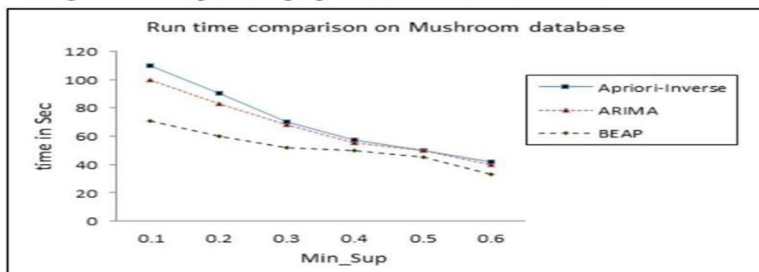


Figure 6: Comparison graph for run time on mushroom database

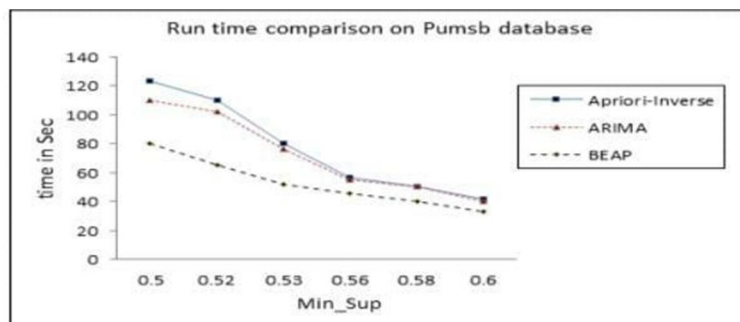


Figure 7: Comparison graph for run time on pumsb database

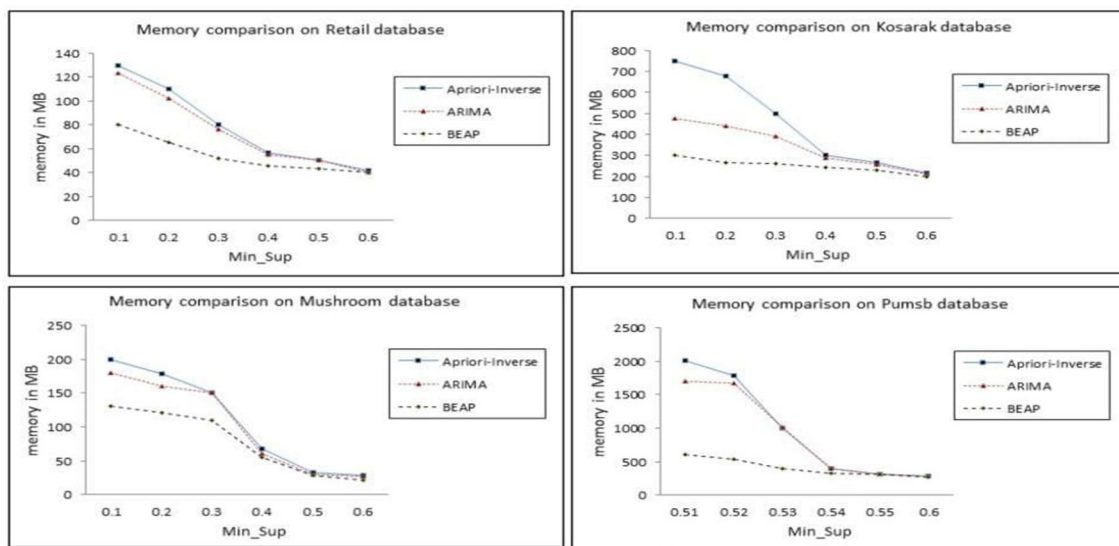


Figure 8: Memory consumption comparison graphs

From Figures 4, 5, 6, 7, it is clear that the proposed BEAP algorithm outperforms other existing algorithms in terms of running time by a good margin. The main reason for the performance of the proposed algorithm is that it only considered rare items and ignored other items, which significantly saved the execution time and outperformed the other two algorithms. The next section shows the memory consumption.

From Figure 8, it is clear that Apriori-inverse performed very poorly on a sparse database like Kosarak and min_sup is above 0.5, all algorithms performed equally well. When the algorithms are implemented on dense databases, the proposed BEAP algorithm outperforms the other two algorithms by a large margin. As shown in Figure 8, it is noted that the proposed BEAP algorithm consumed the least memory for all minimum support threshold values.

IX. CONCLUSION

In this paper a new algorithm BEAP is proposed to detect rare items and the proposed BEAP algorithm is a binary representation and below is an approach to detect rare items with minimal operation time and memory usage. The experimental results show that the proposed BEAP algorithm outperforms the other two algorithms in terms of speed and memory footprint.

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