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Implementation of Pipelined Multi_Precision (1, 2 and 4) Floating-point Arithmetic Operations

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Abstract: The implementation of pipelined multi-precision-based arithmetic operations are carried out. In the existing system, the floating-point operation is based on single precision and is implemented on a divider. The proposed design has been implemented using single, double and quad precision using the universal piece-wise linear (PWL) approximation method and a modified Goldschmidt algorithm. The proposed design performs addition, subtraction, multiplication, and division using the universal PWL method to reduce maximum error. Small multipliers are used in the modified Goldschmidt algorithm. The pipelining process has been used in order to improve the speed of execution and accuracy. This pipelined architecture is described in Verilog and timing performance is measured with Xilinx timing analyser.

Keywords: Pipelined architecture, precision floating Point numbers

I. INTRODUCTION

In the view of the fact that complexity of the algorithms for the floating-point operations are very hard to implement on FPGA. The computations for floating point operations involve large dynamic range, but the resources way required for these operations is high compared with the integer operations. We have unsigned/signed multiplier for multiplication of binary numbers, for multiplication of floating-point numbers floating point division is used. Floating point numbers are one possible of representing real numbers in binary format; The IEEE 754 standard presents two different floating-point formats, Binary interchange format and Decimal interchange format. Dividing floating point numbers is a critical requirement for DSP applications involving large dynamic range. The IEEE Standard for Floating-Point Arithmetic (IEEE 754) is a technical standard for floating-point computation which was established in 1985 by the Institute of Electrical and Electronics Engineers (IEEE IEEE Standard 754 floating point is the most common representation today for real numbers on computers, including Intel-based PC's, Macs, and most Unix platforms.

There are several ways to represent floating point number but IEEE 754 is the most efficient in most cases. IEEE 754 has 3 basic components: The Sign of Mantissa: This is as simple as the name. 0 represents a positive number while 1 represents a negative number. The Biased exponent: The exponent field needs to represent both positive and negative exponents. A bias is added to the actual exponent in order to get the stored exponent. The Normalized Mantissa: The mantissa is part of a number in scientific notation or a floating-point number, consisting of its significant digits. Here we have only 2 digits, i.e., 0 and 1. So a normalized mantissa is one with only one 1 to the left of the decimal. IEEE 754 numbers are divided into Three based on the above three components: single precision and double precision, Quad Precision.

II. METHODOLOGY

MODIFIED GOLDSCHMIDT ALGORITHM is used to find precision values. Hence in Goldschmidt methods, the value of q becomes closer to the accurate value of the result of Y/X with an increasing number of iterations. From the iterative equations, three advantages of the modified Goldschmidt algorithm can be identified, as follows.

- 1) The effective word length of R is reduced. The smaller multiplication and addition units meet the requirements for the iteration over R .
- 2) Complementation is avoided by removing subtraction from the recursive equations.
- 3) The multiplication is replaced by the square operation. the proposed method for multi precision arithmetic operation classified into different modules, namely

XOR Module: XOR is the digital logic gate which provides logic 0 or logic1 as the output. Sign bit performs Xor operation.

Goldschmidt block: The precision values are calculated using modified Goldschmidt algorithm.

Exponent bias addition: The resultant exponent value will be subtracted by the bias value to get the normalized value. The bias value varies based on the precision type.

Floating Point Adder and Subtractor: Addition and subtraction of floating-point numbers is carried out using IEEE 754 standard floating point Addition Algorithm. The resultant sign value decides the arithmetic operation.

Floating Point Multiplier: Multiplication of floating-point numbers is performed by addition of exponents and multiplication of the mantissas.

Floating Point Divider: Division of IEEE 754 Floating point numbers is done by dividing the mantissas and subtracting the exponents.

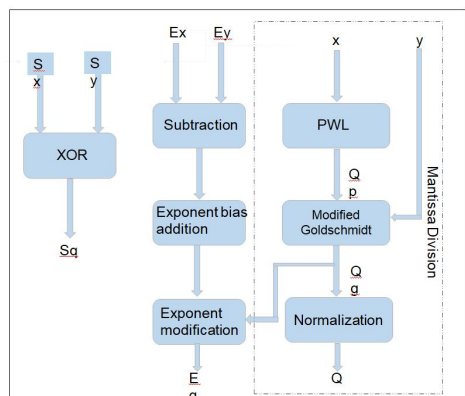


Fig 1: Block Diagram Of Floating-Point Divider

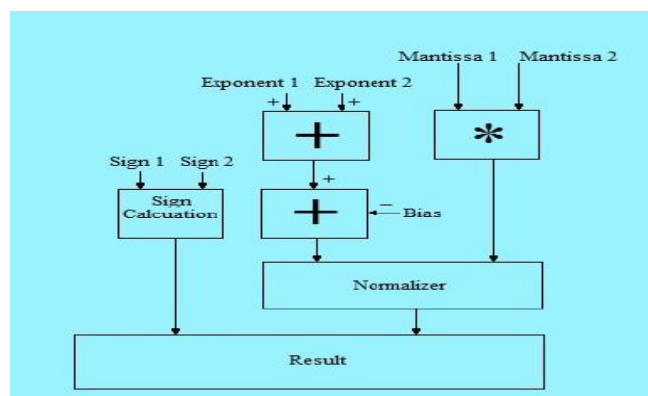


Fig 2: Block Diagram Of Floating-Point Multiplier

III. IMPLEMENTATION

The proposed method for precision floating-point numbers is based on a universal Piece wise linear algorithm method and a modified Goldschmidt algorithm. A state-of-the-art universal PWL method is used to implement the initial approximation of the reciprocal. In the modified Goldschmidt algorithm, full precision multipliers are replaced with smaller multipliers. The design is going to describe a single precision floating point divider for better Area & timing performance. Xilinx Timing Analyzer is used to reduce the power consumption and to increase the speed of execution by implementing certain algorithm for dividing two floating point numbers. By the use of pipelining process this divider is very good at speed and accuracy compared with the previous Dividers. The pipelined architecture is described in Verilog and is implemented on Xilinx Spartan 6 FPGA. Timing performance is measured with Xilinx Timing Analyzer.

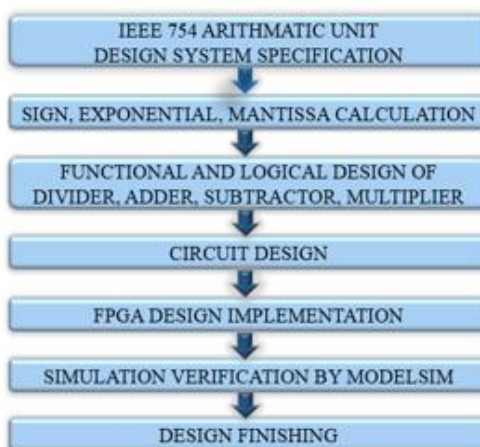


FIG 3:DESIGN FLOW

IV. RESULTS

outputs for the proposed model are explained using an example. Consider 2 inputs namely $x=2.1$ and $y=2.0$. double precision floating point subtraction operation is performed with 2 inputs x & y . Thus, the simulated results in Xilinx are shown in the below fig 4. Hence the results are verified using FPGA spartan 6. obtained result in FPGA spartan 6 is shown in the below fig 5. Table 1 describes the total area and delay utilised in the proposed method, where area and delay are compared with existing model which provides a better performance in terms of area and delay.

Floating_Point_Sub_Double(A	404b999999999999	404b999999999999
Floating_Point_Sub_Double(B	402675c28f5c28f5	402675c28f5c28f5
Floating_Point_Sub_Double(S	40461c28f5c28f5d	40461c28f5c28f5d
Floating_Point_Sub_Double(Bs	1100000000100110	1100000000100110110101110001010001111010111000010100011110110
Floating_Point_Sub_Double(MM(A	0100000000100110	010000000010011011001100110011001100110011001100110011010
Floating_Point_Sub_Double(MM(B	1100000000100110	11000000001001101101011100001010001111010111000010100011110110
Floating_Point_Sub_Double(MM(S	0100000000100110	01000000001001101100001010001111010111000010100011110110
Floating_Point_Sub_Double(MM(Sign_A	S0	
Floating_Point_Sub_Double(MM(Sign_B	S1	
Floating_Point_Sub_Double(MM(Mantissa_A	1011101110011001	10111011100110011001100110011001100110011001100110011010
Floating_Point_Sub_Double(MM(Mantissa_B	0110011101011100	0110011101011100000101000111010111000010100011110110
Floating_Point_Sub_Double(MM(Exponent...	100000000100	100000000100
Floating_Point_Sub_Double(MM(Exponent...	100000000100	100000000100
Floating_Point_Sub_Double(MM(N_E	100000000100	100000000100
Floating_Point_Sub_Double(MM(N_M	0110000111000010	011000011100001010000111010110000101000111101101
Floating_Point_Sub_Double(MM(C	1100000001101111	1100000001101111110111011001110111110111011111101111110

Fig4: simulation of single precision adder



Fig5: snapshot of FPGA

Considering the inputs $x=42.3$ and $y=12.4$ single precision floating point division operation is performed. Hence the result are simulated using Xilinx and are verified using FPGA spartan 6. Thus, the obtained outputs are shown below in fig6 and fig7.

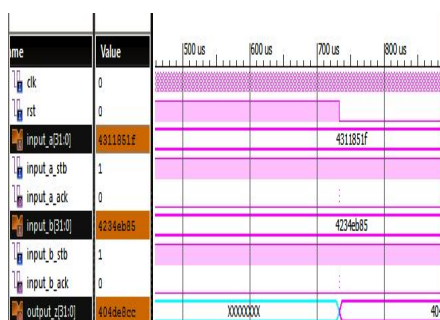


Fig6:simulation of single precision divider

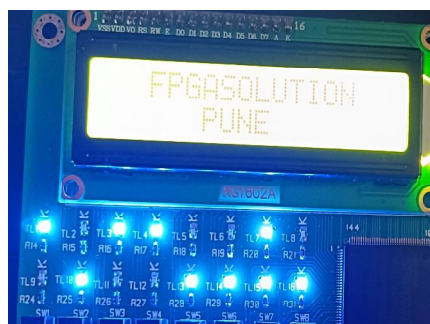


Fig7: snapshot of FPGA

DEVICE NAME	AREA		DELAY
VERTEX 4	LUT	SLICES	OVER DELAY
Single Precision Floating Point Divider	5246	2880	45.6ns
Propose Single Precision Floating Point Divider	862	455	4.692ns

table 1: comparison of synthesized report

V. CONCLUSIONS

The implementation of pipelined multi precision floating point arithmetic operations is executed successfully. The time taken to complete pipelining techniques with delay is noted which is in the order of nano seconds. In comparison to the reference mode, the proposed system obtains better result in terms of area and delay.

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