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32-Bit MIPS RISC Processor

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Abstract: This paper concerned with the design and implementation of a 32 - bit Reduced Instruction Set Computer (RISC) processor on a Field Programmable Gate Arrays (FPGAs). The processor has been designed with VHDL, synthesized using Xilinx ISE 9.1i Webpack, simulated using ModelSim simulator, and then implemented on Xilinx Spartan 2E FPGA that has 143 available Input/Output pins and 50MHz clock oscillator. The results for the different parts of the processor are presented in the form of test bench wave form and the architecture of the system is demonstrated.

Keywords: FPGA, CPU, VHDL, RISC, Processor.

I. INTRODUCTION

Computer Engineering and Computer Design are very much concerned with the cost and performance of components in the implementation domain. Reduced Instruction Set Computer (RISC) focuses on reducing the number and complexity of instructions in the machine [1, 2]. Compared to ASIC design Field Programmable Gate Arrays (FPGAs) are growing fast with cost reduction [3]. The design of a low cost 32 - bit RISC Processor has been described using VHDL, and by using Xilinx FPGA some components have been implemented and tested [4, 5, 6, and 7]. Spartan2E development board, DIO1, and DIO2 extension boards from Digilent have been used for the hardware implementation. For synthesis and simulation the Web pack from Xilinx and ModelSim have been used. The text in this article is organized as follows; The introduction is given in section I; Section II is describe about MIPS processor; Section III is talking about the system architecture; The design of the Control Unit is given in section IV ; Insection V we define the main structure of the Data Path; The design of the ROM is given in section VI; Sections VII will presents the simulation results for the different parts of the processor; The conclusion and future work will be given at the end in section VIII.

II. THE MIPS PROCESSOR

The MIPS direction set engineering (ISA) is a RISC based chip design that was produced by MIPS Computer Systems Inc. in the mid 1980s. MIPS are presently an industry standard and the execution pioneer inside the installed business. Their plans can be found in Canon advanced cameras, Windows CE gadgets, Cisco Routers, Sony Play Station 2 diversion reassures, and numerous more items utilized as a part of our regular day to day existences. By the late 1990s it was assessed that one in three of all RISC chips created was a MIPS-based plan. Engineering of MIPS RISC chip incorporates, settle length straight forward decoded guideline organize, memory gets to constrained to load and store directions, hardwired control unit, a vast universally useful enlist record, and all operations are done inside the registers of the microchip. Because of these outline qualities, PC engineering courses in college and specialized schools the world over frequently concentrate the MIPS design.

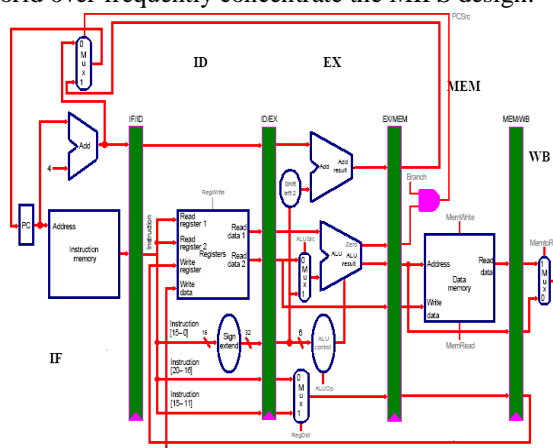


Fig1: The MIPS Processor

The MIPS single-cycle processor plays out the errands of guideline bring, direction disentangle, execution, memory get to and compose back across the board clock cycle. To begin with the PC esteem is utilized as a deliver to list the guideline memory which supplies a 32-bit estimation of the following direction to beexecuted. This direction is then separated into the distinctive fields. The guidelines opcode field bits [31-26] are sent to a control unit to decide the kind of direction to execute. The sort of guideline then figures out which control signs are to be stated and what work the ALU is to perform, along these lines disentangling thedirection. The enroll record peruses in the asked for locations and yields the information values contained in these registers. These information qualities can then be worked on by the ALU whose operation is controlled by the control unit to either register a memory address (e.g. load or store), figure a math result (e.g. include, or potentially slt), or play out a think about (e.g. branch). In the event that the direction decoded is number juggling, the ALU result must be composed to an enlist. On the off chance that the guideline decoded is a heap or a store, the ALU result is then used to address the information memory. The last stride composes the ALU result or memory esteem back to the enroll record. The MIPS usage as with all processors comprises of two principle sorts of rationale components: combinational and successivecomponents.

III. SYSTEMARCHITECTURE

Figure.2 shows the three components of the RISC processor presented in this paper. The three components are the Control Unit (CU), the Data Path, and the ROM. The Central Processing Unit (CPU) has 17 instructions. The design of the processor will be describe in the followingsections.

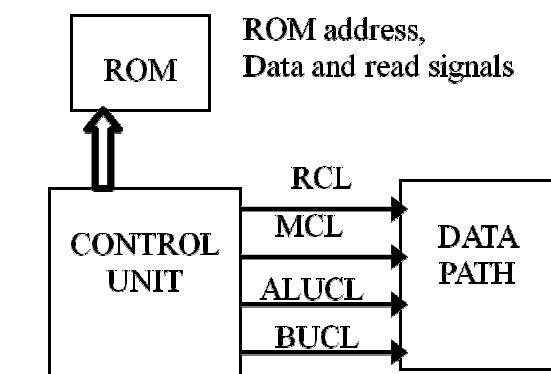


Figure (2) SystemArchitecture

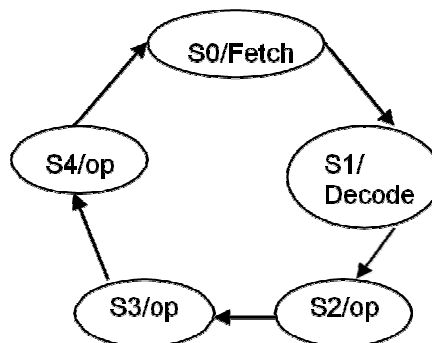


Figure (3) State diagram of CUDesign

IV. DESIGN OF THE CONTROL UNIT

By using FSM (Finite State Machine) the control unit is designed and it allows each state to run at one clock cycle, the first state is the reset which initialize the CPU internal registers and variables. The machine goes to the reset state by enabling the reset signal for a certain number of clocks. The instruction fetching and decoding states are followed by the reset state. The instruction state is for reading instruction data from the ROM and the decoding state is decodes the parts of the instruction. Depending on the instruction the decoding state will also select the next state, since every instruction has its own set of states, based on the given instruction the control unit will jump to the correct state. After completion of all states of a running instruction, the last one will allow us to process the next instruction in the program so it returns to the fetch state. The state diagram of the control unit as shown in Figure.3.

V. DESIGN OF THE DATA PATH

The Data Path consists of subunits that are necessary for performing all of arithmetic and logic operations. A Datapath is a hardware that performs data processing operations [8, 9, 10, and 11]. It is one of two types of modules used to represent a digital system, the other being a control unit. The Datapath model we designed consists of the units necessary to perform all the operations on the data selected by the control unit. Figure.3 shows the components include a Register File, Arithmetic/Logic Unit, Memory Interface and Branching Unit. The Register File holds the table of the 32 general purpose registers available to the CPU, it has two output ports (output1, output2) and one input port, and also it has a 16 bit bus connected directly to the Control Unit to pass immediatedata.

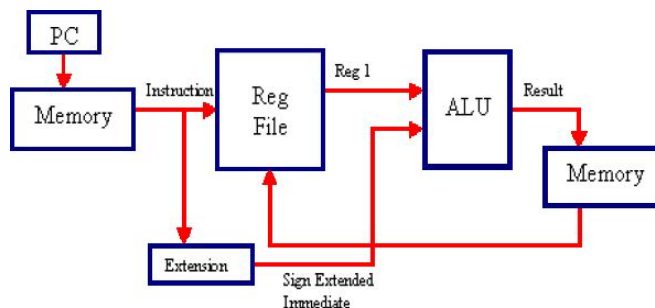


Fig 4: Load word Data Path

The *ALU* design consists of two input ports and one output port which mainly performs operations on two operands. It has a design similar to the control unit based on a code given by the *ALUCL* it selects an operation. The *Memory Interface* was designed to accommodate simple load/store operations with the 16x32 memory. The effective address is calculated by adding the content of the address register and the immediate data. The Branch Unit calculates a given condition by the control unit and raises a branch flag whether the condition is met or not, and if the flag is raised, it sends the branch address back to the control unit in order to replace the program counter. The control lines coming from the control unit operate all the units in the data path. The path starts from the register file that has two output ports which are connected to all the other units, after that the processing is done by one of the other units then finally returned back to the register files input port using the multiplexer. The signals used in the data path are forwarded from the control unit to each subcomponent as needed.

MCL

BUCL

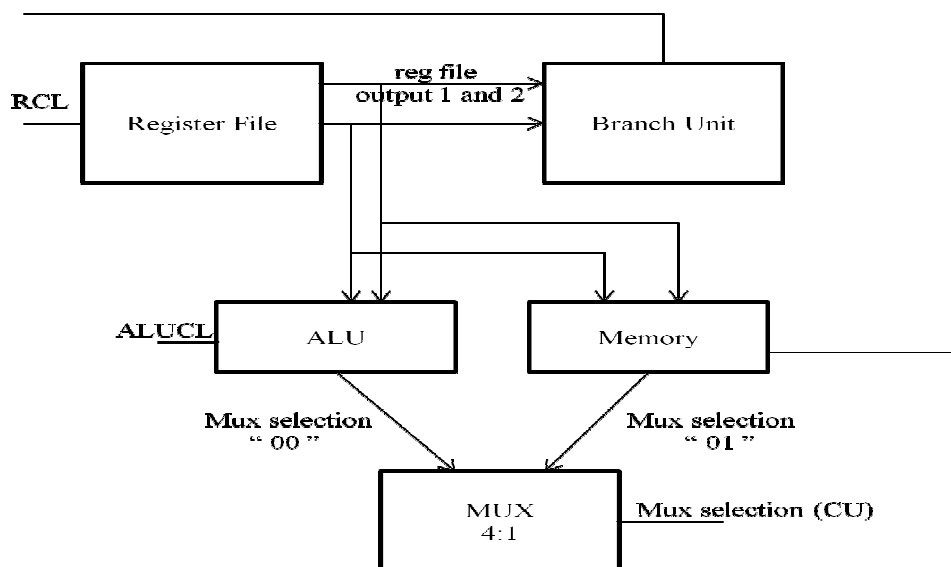


Figure (5) the Sub-components of the Data Path

VI. DESIGN OF THE ROM

The CPU has a built in ROM which enables us to program simple code and execute it. It is a basic 16x32 ROM and it is 32bit aligned. The List of signals in the ROM is:

- A. address: address sent by the control unit
- B. data_out: the data that is contained the given address
- C. read: the signal to enable reading from the ROM
- D. ready: the signal to indicate when the ROM is ready for reading
- E. clk: the main clock signal
- F. reset: the initial reset signal

VII.SIMULATIONRESULTS

Depictions of each snapshot of the application while running is given underneath:

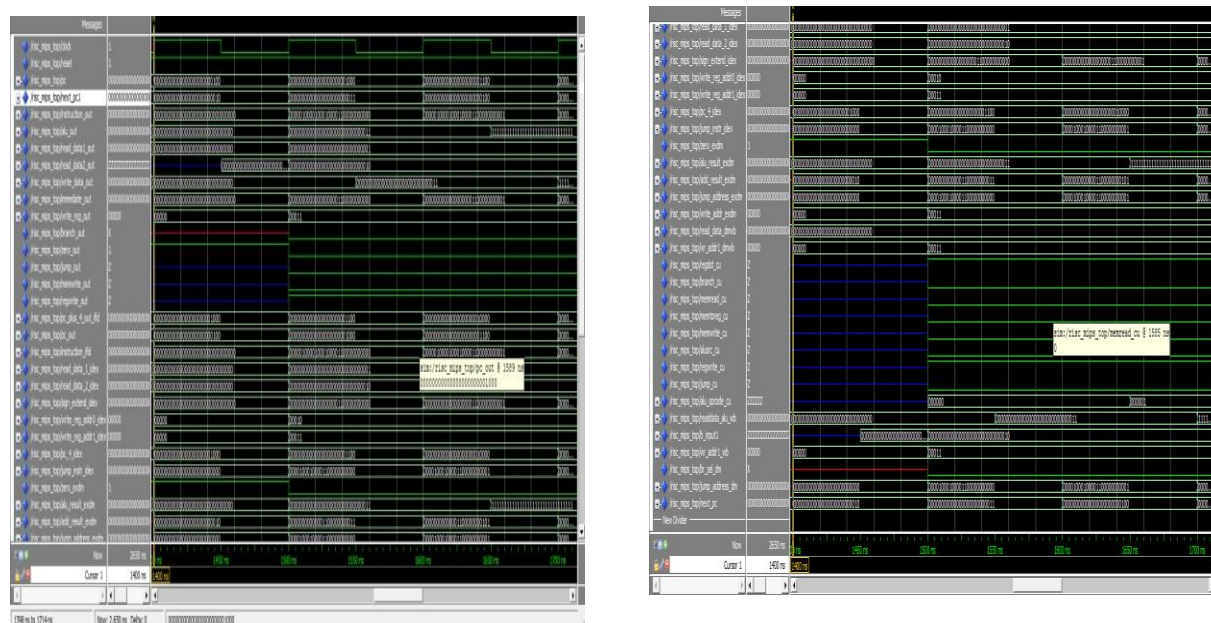


Figure (6) Simulation results

VIII. CONCLUSION AND FUTUREWORK

32-bit RISC Process has been design and implemented in hardware on Xilinx Spartan 2E FPGA. The design has been achieved using VHDL and simulated with ModelSim. The hardware part has been used by the Digilent Spartan 2E development board. Most of the goals were achieved and simulation shows that the processor is working perfectly, but the Spartan 2E FPGA was not sufficient for implementing the whole design into a real hardware, since the total available logic gate in Spartan 2E is 200K Logic Gate, which was not enough for implementing the whole processor, but parts of the processor have been implemented and test in a real hardware. In future scope increasing the number of instructions and make a pipelined design with less clock cycles per instruction and more improvement can be added in the future work.

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