



iJRASET

International Journal For Research in
Applied Science and Engineering Technology



INTERNATIONAL JOURNAL FOR RESEARCH

IN APPLIED SCIENCE & ENGINEERING TECHNOLOGY

Volume: 3

Issue: V

Month of publication: May 2015

DOI:

www.ijraset.com

Call: ☎ 08813907089

E-mail ID: ijraset@gmail.com

A Load Balancing Model Based on Cloud Partitioning for the Public Cloud

Sakshi A. Rayake, Asst. Prof. P. Chapke

Abstract: *Load balancing in the cloud computing environment has an important impact on the performance. Good load balancing makes cloud computing more efficient and improves user satisfaction. This article introduces a better load balance model for the public cloud based on the cloud partitioning concept with a switch mechanism to choose different strategies for different situations. The algorithm applies the game theory to the load balancing strategy to improve the efficiency in the public cloud environment.*

I. INTRODUCTION

Cloud computing is an attracting technology in the field of computer science. In Gartner's report [1], it says that the cloud will bring changes to the IT industry. The cloud is changing our life by providing users with new types of services. Users get service from a cloud without paying attention to the details [2]. NIST gave a definition of cloud computing as a model for enabling ubiquitous, convenient, on-demand network access to a shared pool of configurable computing resources (e.g., networks, servers, storage, applications, and services) that can be rapidly provisioned and released with minimal management effort or service provider interaction [3]. More and more people pay attention to cloud computing [4, 5]. Cloud computing is efficient and scalable but maintaining the stability of processing so many jobs in the cloud computing environment is a very complex problem with load balancing receiving much attention for researchers. Since the job arrival pattern is not predictable and the capacities of each node in the cloud differ, for load balancing problem, workload control is crucial to improve system performance and maintain stability. Load balancing schemes depending on whether the system dynamics are important can be either static and dynamic [6]. Static schemes do not use the system information and are less complex while dynamic schemes will bring additional costs for the system but can change as the system status changes. A dynamic scheme is used here for its flexibility. The model has a main controller and balancers to gather and analyze the information. Thus, the dynamic control has little influence on the other working nodes. The system status then provides a basis for choosing the right load balancing strategy. The load balancing model given in this article is aimed at the public cloud which has numerous nodes with distributed computing resources in many different geographic locations. Thus, this model divides the public cloud into several cloud partitions. When the environment is very large and complex, these divisions simplify the load balancing. The cloud has a main controller that chooses the suitable partitions for arriving jobs while the balancer for each cloud partition chooses the best load balancing strategy.

II. REALATED WORK

There have been many studies of load balancing for the cloud environment. Load balancing in cloud computing was described in a white paper written by Adler[7] who introduced the tools and techniques commonly used for load balancing in the cloud. However, load balancing in the cloud is still a new problem that needs new architectures to adapt to many changes. Chaczko et al.[8] described the role that load balancing plays in improving the performance and maintaining stability. There are many load balancing algorithms, such as Round Robin, Equally Spread Current Execution Algorithm, and Ant Colony algorithm. Nishant et al.[9] used the ant colony optimization method in nodes load balancing. Randles et al.[10] gave a compared analysis of some algorithms in cloud computing by checking the performance time and cost. They concluded that the ESCE algorithm and throttled algorithm are better than the Round Robin algorithm. Some of the classical loads balancing methods are similar to the allocation method in the operating system, for example, the Round Robin algorithm and the First Come First Served (FCFS) rules. The Round Robin algorithm is used here because it is fairly simple.

III. SYSTEM MODEL

There are several cloud computing categories with this work focused on a public cloud. A public cloud is based on the standard cloud computing model, with service provided by a service provider [11]. A large public cloud will include many nodes and the nodes in different geographical locations. Cloud partitioning is used to manage this large cloud. A cloud partition is a subarea of the public cloud with divisions based on the geographic locations. The architecture is shown in Fig.1.

International Journal for Research in Applied Science & Engineering Technology (IJRASET)

The load balancing strategy is based on the cloud partitioning concept. After creating the cloud partitions, the load balancing then starts: when a job arrives at

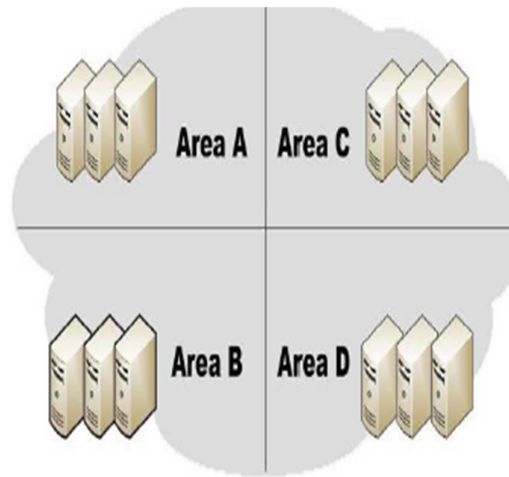


Fig. 1 Typical cloud partitioning.

the system, with the main controller deciding which cloud partition should receive the job. The partition load balancer then decides how to assign the jobs to the nodes. When the load status of a cloud partition is normal, this partitioning can be accomplished locally. If the cloud partition load status is not normal, this job should be transferred to another partition. The whole process is shown in Fig.2.

A. Main controller and balancers

The load balance solution is done by the main controller and the balancers. The main controller first assigns jobs to the suitable cloud partition and then communicates with the balancers in each partition to refresh this status information. Since the main controller deals with information for each partition, smaller data sets will lead to the higher processing rates. The balancers in each partition gather the status information from every node and then choose the right strategy to distribute the jobs. The relationship between the balancers and the main controller is shown in Fig.3.

B. Assigning jobs to the cloud partition

When a job arrives at the public cloud, the first step is to choose the right partition. The cloud partition status can be divided into three types:

- (1) Idle: When the percentage of idle nodes exceeds α change to idle status.

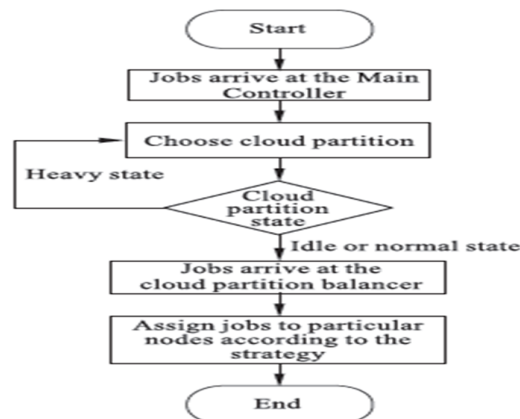


Fig. 2 Job assignment strategy.

International Journal for Research in Applied Science & Engineering Technology (IJRASET)

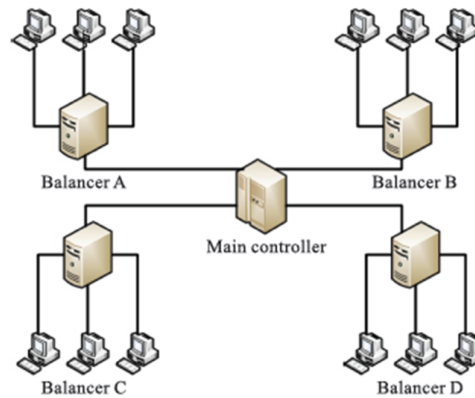


Fig. 3 Relationships between the main controllers, the balancers, and the nodes.

(2) Normal: When the percentage of the normal nodes exceeds β , change to normal load status.

(3) Overload: When the percentage of the overloaded nodes exceeds γ , change to overloaded status.

The parameters α , β and γ are set by the cloud partition balancers. The main controller has to communicate with the balancers frequently to refresh the status information. The main controller then dispatches the jobs using the following strategy:

When job i arrives at the system, the main controller queries the cloud partition where job is located. If this location's status is idle or normal, the job is handled locally. If not, another cloud partition is found that is not overloaded. The algorithm is shown in Algorithm 1.

3.3 Assigning jobs to the nodes in the cloud

Partition The cloud partition balancer gathers load information from every node to evaluate the cloud partition status. This evaluation of each node's load status is very important. The first task is to define the load degree of each nodes.

Algorithm 1 Best Partition Searching

```

begin
while job do
  searchBestPartition (job);
  if partitionState == idle k partitionState == normal then
    Send Job to Partition;
  else
    search for another Partition;
  end if
end while
end

```

The node load degree is related to various static parameters and dynamic parameters. The static parameters include the number of CPU's, the CPU processing speeds, the memory size, etc. Dynamic parameters are the memory utilization ratio, the CPU utilization ratio, the network bandwidth, etc. The load degree is computed from these parameters as below:

Step 1 Define a load parameter set: $F = \{F_1, F_2, \dots, F_m\}$ with each F_i ($1 \leq i \leq m, F_i \in [0, 1]$)

Parameter being either static or dynamic. M represents the total number of the parameters.

Step 2 Compute the load degree as:

$$\text{Load_degree}_{\text{avg}} = \frac{\sum_{i=1}^n \text{Load_degree}(N_i)}{n}$$

The bench mark Load degree_{high} is then set for different situations based on the Load degree_{avg}.

Step 4 Three nodes load status levels are then

Defined as:

International Journal for Research in Applied Science & Engineering Technology (IJRASET)

- Idle When

Load degree.N = 0.

There is no job being processed by this node so the status is charged to Idle.

- Normal For

$0 < \text{Load degree}(N) \leq \text{Load degreehigh}$;

The node is normal and it can process other jobs.

- Overloaded When

Load degreehigh \leq Load degree(N).

The node is not available and can not receive jobs until it returns to the normal.

The load degree results are input into the Load Status

Tables created by the cloud partition balancers. Each balancer has a Load Status Table and refreshes it each fixed period T.

The table is then used by the balancers to calculate the partition status. Each partition status has a different load balancing solution. When a job arrives at a cloud partition, the balancer assigns the job to the nodes based on its current load strategy.

This strategy is changed by the balancers as the cloud partition status changes.

IV. CLOUD PARTITION LOAD BALANCING STRATEGY

A. Motivation

Good load balance will improve the performance of the entire cloud. However, there is no common method that can adapt to all possible different situations. Various methods have been developed in improving existing solutions to resolve new problems. Each particular method has advantage in a particular area but not in all situations. Therefore, the current model integrates several methods and switches between the load balance method based on the system status. A relatively simple method can be used for the partition idle state with a more complex method for the normal state. The load balancers then switch methods as the status changes. Here, the idle status uses an improved Round Robin algorithm while the normal status uses a game theory based load balancing strategy

B. Load balance strategy for the idle status

When the cloud partition is idle, many computing resources are available and relatively few jobs are arriving. In this situation, this cloud partition has the ability to process jobs as quickly as possible so a simple load balancing method can be used. There are many simple load balance algorithm methods such as the Random algorithm, the Weight Round Robin, and the Dynamic Round Robin[12]. The Round Robin algorithm is used here for its simplicity. The Round Robin algorithm is one of the simplest load balancing algorithms, which passes each new request to the next server in the queue. The algorithm does not record the status of each connection so it has no status information. In the regular Round Robin algorithm, every node has an equal opportunity to be chosen. However, in a public cloud, the configuration and the performance of each node will be not the same. Thus, this method may overload some nodes. Thus, an improved Round Robin algorithm is used, which called "Round Robin based on the load degree evaluation". The algorithm is still fairly simple. Before the Round Robin step, the nodes in the load balancing table are ordered based on the load degree from the lowest to the highest. The system builds a circular queue and walks through the queue again and again. Jobs will then be assigned to nodes with low load degrees. The node order will be changed when the balancer refreshes the Load Status Table.

However, there may be read and write inconsistency at the refresh period T. When the balance table is refreshed, at this moment, if a job arrives at the cloud partition, it will bring the inconsistent problem. The system status will have changed but the information will still be old. This may lead to an erroneous load strategy choice and an erroneous nodes order. To resolve this problem, two Load Status Tables should be created as: Load Status Table 1 and Load Status Table 2. A flag is also assigned to each table to indicate Read or Write. When the flag = "Read", then the Round Robin based on the load degree evaluation algorithm is using this table. When the flag = "Write", the table is being refreshed, new information is written into this table. Thus, at each moment, one table gives the correct node locations in the queue for the improved Round Robin algorithm, while the other is being prepared with the updated information. Once the data is refreshed, the table flag is changed to "Read" and the other table's flag is changed to "Write". The two tables then alternate to solve the inconsistency. The process is shown in Fig.4.

C. Load balancing strategy for the normal status

When the cloud partition is normal, jobs are arriving much faster than in the idle state and the situation is far more complex, so a different strategy is used for the load balancing. Each user wants his jobs completed in the shortest time, so the public

International Journal for Research in Applied Science & Engineering Technology (IJRASET)

cloud needs a method that can complete the jobs of all users with reasonable response time. The load balancing in the cloud computing environment can be viewed as a game. Game theory has non-cooperative games and cooperative games. In cooperative games, the decision makers eventually come to an agreement which is called a binding agreement. Each decision maker decides by comparing notes with each others. In non-cooperative games, each decision maker makes decisions only for his own benefit. The system then reaches the Nash equilibrium, where each decision maker makes the optimized decision. The Nash equilibrium is when each player in the game has chosen a strategy and no player can benefit by changing his or her strategy while the other players strategies remain unchanged. There have been many studies in using game theory for the load balancing. Grosu et al.[14] proposed a load balancing strategy based on game theory for the distributed systems as a non-cooperative game using the distributed structure. They compared this algorithm with other traditional methods to show that their algorithm was less complexity with better performance. Aote and Kharat[15] gave a dynamic load balancing model based on game theory. This model is related on the dynamic load status of the system with the users being the decision makers in a non-cooperative game. Since the grid computing and cloud computing environments are also distributed system, these algorithms can also be used in grid computing and cloud computing environments. Previous studies have shown that the load balancing strategy for a cloud partition in the normal load status can be viewed as a no cooperative game, as described here. The players in the game are the nodes and the jobs. Suppose there are n nodes in the current cloud partition with N jobs arriving, and then define the following parameters:

μ_i : Processing ability of each node, $i = 1 \dots n$.

Φ_j : Time spending of each job.

Appropriate value of s_{ji} . The current model uses the method of Grosu et al.[14] called “the best reply” to calculate s_{ji} of each node, with a greedy algorithm then used to calculate s_{ji} for all nodes. This procedure gives the Nash equilibrium to minimize the response time of each job. The strategy then changes as the node’s statuses change.\

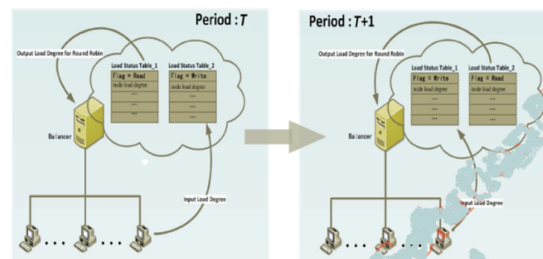


Fig. 4 The solution of inconsistency problem.

V. ACKNOWLEDGEMENT

We would like to thank the editors and anonymous reviewers for their valuable comments and helpful suggestions.

REFERENCES

- [1] R. Hunter, The why of cloud, <http://www.gartner.com/DisplayDocument?doc cd=226469&ref= g noreg>, 2012.
- [2] M. D. Dikaiakos, D. Katsaros, P. Mehra, G. Pallis, and A. Vakali, Cloud computing: Distributed internet computing for IT and scientific research, *Internet Computing*, vol.13, no.5, pp.10-13, Sept.-Oct. 2009.
- [3] P. Mell and T. Grance, The NIST definition of cloud computing, <http://csrc.nist.gov/publications/nistpubs/800-145/SP800-145.pdf>, 2012.
- [4] Microsoft Academic Research, Cloud computing, <http://libra.msra.cn/Keyword/6051/cloud-computing?query= cloud%20computing>, 2012.
- [5] Google Trends, Cloud computing, <http://www.google.com/trends/explore#q=cloud%20computing>, 2012
- [6] N. G. Shivaratri, P. Krueger, and M. Singhal, Load distributing for locally distributed systems, *Computer*, vol. 25, no. 12, pp. 33-44, Dec. 1992.
- [8] Z. Chaczko, V. Mahadevan, S. Aslanzadeh, and C. Mcdermid, Availability and load balancing in cloud computing, presented at the 2011 International Conference on Computer and Software Modeling, Singapore, 2011.
- [9] K. Nishant, P. Sharma, V. Krishna, C. Gupta, K. P. Singh, N. Nitin, and R. Rastogi, Load balancing of nodes in cloud using ant colony optimization, in *Proc. 14th International Conference on Computer Modelling and Simulation (UKSim)*, Cambridgeshire, United Kingdom, Mar. 2012, pp. 28-30.
- [10] M. Randles, D. Lamb, and A. Taleb-Bendiab, A comparative study into distributed load balancing algorithms for cloud computing, in *Proc. IEEE 24th International Conference on Advanced Information Networking and Applications*, Perth, Australia, 2010, pp. 551-556.

International Journal for Research in Applied Science & Engineering Technology (IJRASET)

- [11] A. Rouse, Public cloud, <http://searchcloudcomputing.techtarget.com/definition/public-cloud>, 2012.
- [12] D. MacVittie, Intro to load balancing for developers — The algorithms, <https://devcentral.f5.com/blogs/us/intro-to-load-balancing-for-developers-ndash-the-algorithms>, 2012.
- [13] S. Penmatsa and A. T. Chronopoulos, Game-theoretic static load balancing for distributed systems, *Journal of Parallel and Distributed Computing*, vol. 71, no. 4, pp. 537-555, Apr. 2011.
- [14] D. Grosu, A. T. Chronopoulos, and M. Y. Leung, Load balancing in distributed systems: An approach using cooperative games, in *Proc. 16th IEEE Intl. Parallel and Distributed Processing Symp.*, Florida, USA, Apr. 2002, pp. 52-61.
- [15] S. Aote and M. U. Kharat, A game-theoretic model for dynamic load balancing in distributed systems, in *Proc. The International Conference on Advances in Computing, Communication and Control (ICAC3 '09)*, New York, USA, 2009, pp. 235-238



10.22214/IJRASET



45.98



IMPACT FACTOR:
7.129



IMPACT FACTOR:
7.429



INTERNATIONAL JOURNAL FOR RESEARCH

IN APPLIED SCIENCE & ENGINEERING TECHNOLOGY

Call : 08813907089  (24*7 Support on Whatsapp)